

Thm (VCG): every VCG mechanism is incentive

truthful:  $a = f(v_i, v_{-i})$   
 $v_i(a) + \sum_{j \neq i} v_j(a) - h_i(v_{-i})$

lies:  $a' = f(v_i', v_{-i})$   
 $v_i(a') + \sum_{j \neq i} v_j(a') - h_i(v_{-i})$

But since  $a = f(v_i, v_{-i})$  maximizes social welfare over all alternatives,

$$v_i(a) + \sum_{j \neq i} v_j(a) \geq v_i(a') + \sum_{j \neq i} v_j(a')$$

and thus the same inequality

holds when we subtract term  $h_i(v_{-i})$  from both sides.

Ex. in Auction of a single item, finding the player with highest value is exactly equivalent to maximizing  $\sum_i v_i(i)$  since only a single player gets non-zero value. The payment is Clarke pivot payment below

If all  $h_i = 0$ , then though the mechanism is simple but the mechanism pays money.

Def: The choice  $h_i(v_{-i}) = \max_{b \in A} \sum_{j \neq i} v_j(b)$

is called the Clarke pivot payment. Under this rule the

payment of player  $i$  is  $p_i(v_1, \dots, v_n) = \max_b \sum_{j \neq i} v_j(b) - \sum_{j \neq i} v_j(a)$   
where  $a = f(v_1, \dots, v_n)$ .

Intuitively,  $i$  pays an amount equal to the total damage that he causes the other players. The difference between the social welfare of the others with and without  $i$ 's participation.

Lim: A VCG with Clarke pivot payments makes the case that no players ~~always get non-negative utility~~ is ever paid.

Pf: Trivial

Clarke pivot rules does not fit many situations where valuations are negative but it is fairly general rule  
payment. negative